

# Secure and Efficient Access to Outsourced Data

Weichao Wang, Zhiwei Li, Rodney Owens, Bharat  
Bhargava

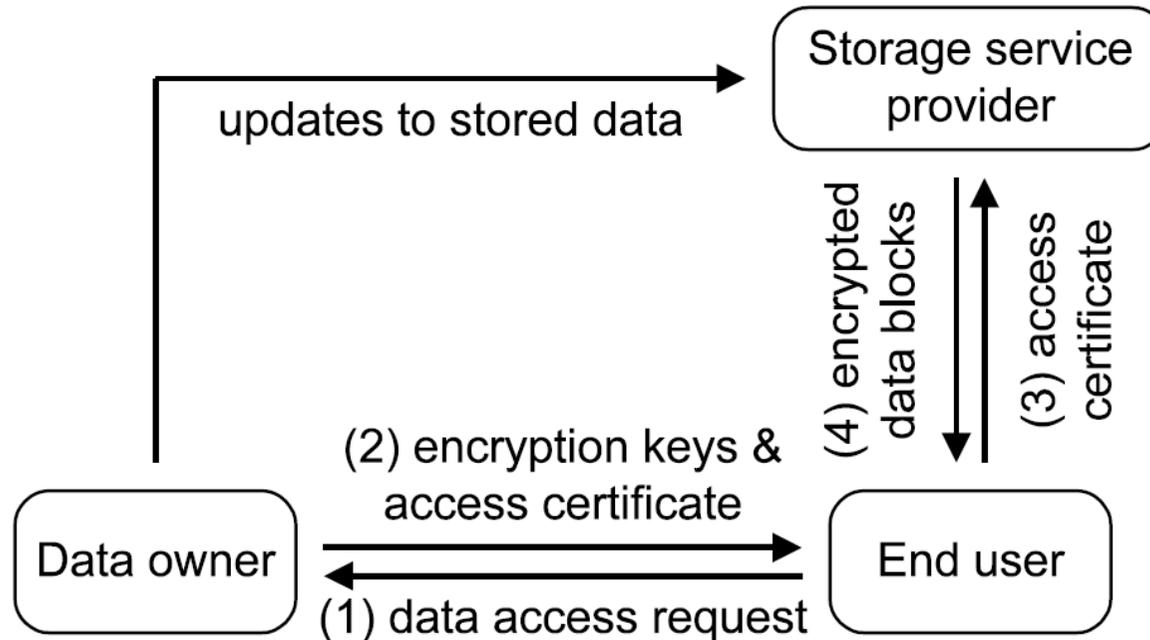
CCSW 2009: The ACM Cloud Computing Security Workshop

# The Problem

- Providing secure and efficient access to outsourced data
  - An important component of cloud computing
  - Foundation for information management and other operations
- the security guidance published by Cloud Security Alliance
  - strong encryption and scalable key management
  - information lifecycle management
  - system availability and performance

# Investigated Environment

- Owner-write-user-read Scenario
  - Data can be updated only by the original owner
  - Users read the information according to access rights
  - Example Application: LHC (Large Hadron Collider)



# The Solution

- Fine grained access control to outsourced data
  - encrypt every data block with a different symmetric key
- Flexible and efficient management
  - adopt the key derivation method to reduce the number of secrets maintained
- Data isolation among end users
  - adopt over-encryption
  - lazy revocation
- Mechanisms to handle dynamics in both user access rights and outsourced data

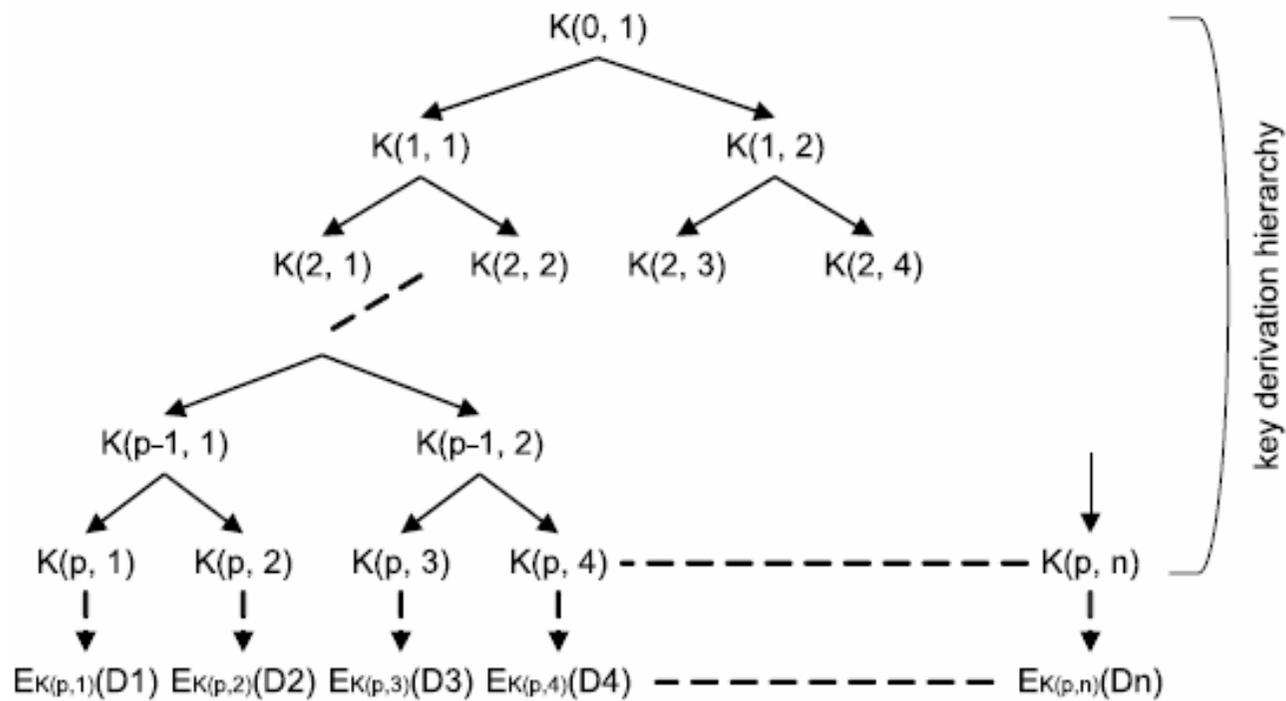
# Fine grained access control

- Encrypt every data block with a different symmetric key
  - Data blocks  $\{D_1, D_2, \dots, D_n\}$
  - Encryption keys  $k_i$  ( $i=1$  to  $n$ )
  
- Worst case
  - Storage overhead linear to  $n$
  - Communication overhead linear to  $n$

# Key-derivation-based data block encryption

- Key derivation method
  - Generate the data block encryption keys through a hierarchy
  - Every key in the hierarchy can be derived by combining its parent node and some public information
  - Calculation of one-way functions

# Key derivation hierarchy



left child of  $K(i, j)$ :  $K(i+1, (2 * j - 1)) = \text{hash}(K(i, j) || (2 * j - 1) || K(i, j))$

right child of  $K(i, j)$ :  $K(i+1, (2 * j)) = \text{hash}(K(i, j) || (2 * j) || K(i, j))$

# Issues of the key hierarchy

- Account for data updates
  - leave some room for the insertion and appending operations
- Only distribute necessary keys
  - we should not disclose keys of the blocks that are temporarily missing
- Impact of users' access rights on the communication overhead
  - organize data blocks with similar access patterns into groups

# Data Access Procedure

1. (End user) sends a data access request to the data owner

$$\mathcal{U} \rightarrow \mathcal{O} : \{ \mathcal{U}, \mathcal{O}, E_{k_{\mathcal{O}\mathcal{U}}}(\mathcal{U}, \mathcal{O}, \text{request index}, \text{data block indexes}, \text{MAC code}) \}$$

2. (Data owner) authenticate the sender, verify the request, and determine the smallest key set

$$\mathcal{O} \rightarrow \mathcal{U} : \{ \mathcal{O}, \mathcal{U}, E_{k_{\mathcal{O}\mathcal{U}}}(\mathcal{O}, \mathcal{U}, \text{request index}, \text{ACM index}, \text{seed for } P(), \mathcal{K}', \text{cert for } \mathcal{S}, \text{MAC code}) \}$$

- $K$
- *ACM index*
- *cert*

$$\{ E_{k_{\mathcal{O}\mathcal{S}}}(\mathcal{U}, \text{request index}, \text{ACM index}, \text{seed}, \text{indexes of data blocks}, \text{MAC code}) \}$$

## Data Access Procedure

3. (End user) sends  $\{U, S, request\ index, cert\}$  to the service provider
4. (Service provider) verify the *cert*, check the user and ACM index, and retrieve data blocks and conduct the over-encryption
5. (End user) receive the data blocks, use seed and  $K'$  to derive keys, and then recover the data

# Over-encryption

- Confidentiality of the outsourced data
  - Prevent revoked users from getting access to out-sourced data through eavesdropping
- $P()$ : a pseudo random bit sequence generator
  - Shared between service provider and end users
- Given a *seed*,  $P()$  can generate a sequence of pseudo random bits
- Procedure
  - Use *seed* and  $P()$  generate a sequence of pseudo random bits
  - Use this bit sequence as one-time pad xor it to the encrypted block

# Dynamics in User Access Rights

- Grant Access Right
  - Change access control matrix
  - Increase the value of *ACM* index
  - Service provider and the end user do not need to change

# Dynamics in User Access Rights

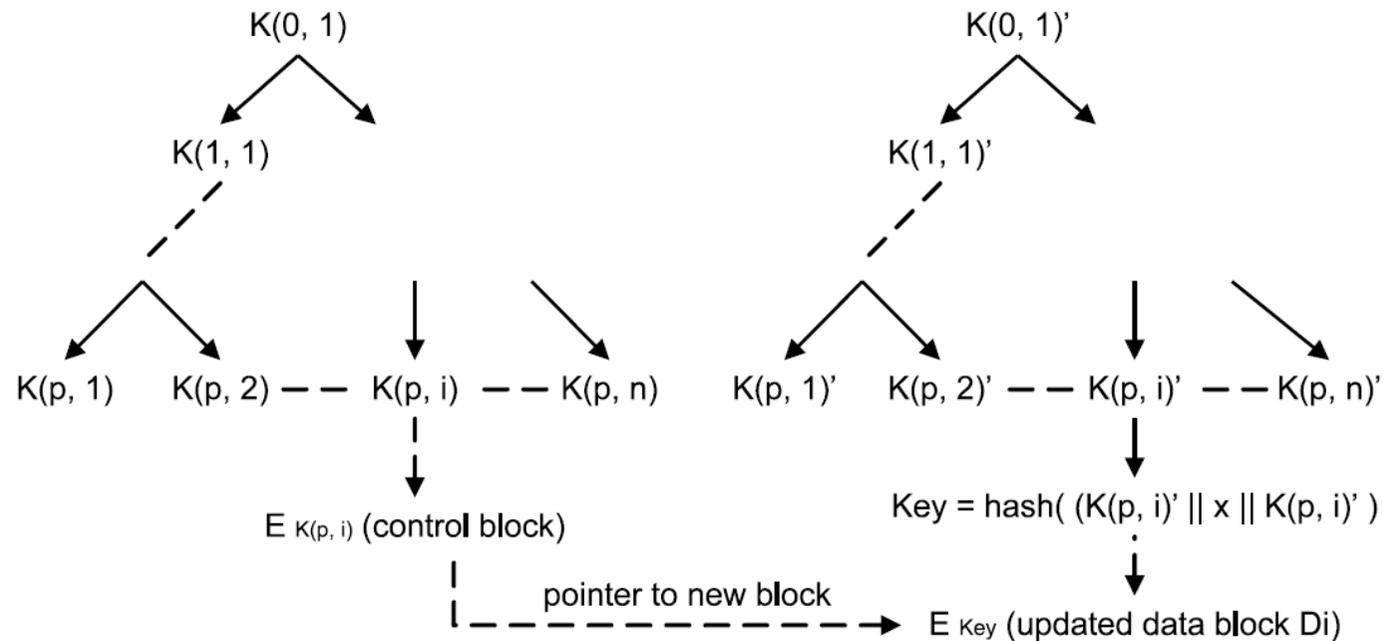
- Revoke Access Right
  - Depends on whether or not the service provider conducts over-encryption
- If service provider conducts over-encryption
  - (Owner) updates the access control matrix and increase the ACM index
  - (Owner) send the new ACM index to the service provider until it receives acknowledgement
- If service provider refuses to conducts over-encryption
  - Adopt the lazy revocation method to prevent end users from reading updated blocks
  - trades re-encryption and data access overhead for a degree of security

# Dynamics in Outsourced Data

- Block Deletion
  - use a special control block to replace
  - label non-existence in the access control matrix
- Block Insertion / Appending
  - locate an unused block index
  - derive the encryption key
  - encrypt the data block
  - store it on the service provider
  - insert new data blocks based on their access patterns

# Dynamics in Outsourced Data

- Block Update



Control block:

- (1). Pointer to the new data block
- (2). Information used to derive the encryption key of  $D_i'$
- (3). Information to verify integrity

# Overhead of the proposed approach

computational overhead (in machine cycle)			
	owner $\mathcal{O}$	server $\mathcal{S}$	user $\mathcal{U}$
key derivation	27M	–	720M
one-time pad generation and over-encryption	–	10G	10G
communication overhead			
	owner $\mathcal{O}$	server $\mathcal{S}$	user $\mathcal{U}$
data blk index #	6KByte	–	10KByte
control blk	–	–	10.5KByte
keys in hierarchy	16KByte	–	–
updated data blk	–	1MByte	–

Outsourced data size: 10 PB

Data block size: 4 KB

Key hierarchy height:  $p = 42$

User retrieve 1GB=250,000 blocks

## Comparison to approach proposed by Atallah et al. (CCS'05)

- Their approach is more generic
- However, our approach
  - has less communication and storage overhead for data retrieval when they have infrequent update operations
  - handles user revocation without impacting service provider (over-encryption, lazy-revocation)

# Conclusion

- Propose a mechanism to achieve secure and efficient access to outsourced data in owner-write-users-read applications.
- Analysis shows that the key derivation procedure based on hash functions will introduce very limited overhead.
- Use over-encryption and/or lazy revocation to prevent revoked users from getting access to updated data blocks.
- We design mechanisms to handle both updates to outsourced data and changes in user access rights.

# Future work

- Design a new scheme for key management for many-write-many-read applications
- Further reduce the number of keys by recognizing the access patterns of data blocks
- Develop a new approach to secure Storage-as-a-Service.